

Bandit View on Continuous Stochastic Optimization

Sébastien Bubeck¹

joint work with Rémi Munos¹ & Gilles Stoltz² & Csaba Szepesvari³

¹ INRIA Lille, SequeL team

² CNRS/ENS/HEC

³ University of Alberta

\mathcal{X} -armed bandit game

Parameters available to the forecaster: the number of rounds n and the set of arms \mathcal{X} .

Parameters unknown to the forecaster: mean-payoff function $f : \mathcal{X} \rightarrow [0, 1]$, reward distributions (over $[0, 1]$) $M(x)$ such that $f(x)$ is the expectation of $M(x)$.

For each round $t = 1, 2, \dots, n$;

- ① The player chooses an arm $X_t \in \mathcal{X}$.
- ② The environment draws the reward Y_t from $M(X_t)$ (and independently from the past given X_t).

Goal: Maximize (in expectation) the cumulative rewards.
Equivalently we want to minimize the cumulative regret

$$R_n = \mathbb{E} \sum_{t=1}^n \left(\max_{x \in \mathcal{X}} f(x) - Y_t \right).$$

\mathcal{X} -armed bandit game

Parameters available to the forecaster: the number of rounds n and the set of arms \mathcal{X} .

Parameters unknown to the forecaster: mean-payoff function $f : \mathcal{X} \rightarrow [0, 1]$, reward distributions (over $[0, 1]$) $M(x)$ such that $f(x)$ is the expectation of $M(x)$.

For each round $t = 1, 2, \dots, n$;

- 1 The player chooses an arm $X_t \in \mathcal{X}$.
- 2 The environment draws the reward Y_t from $M(X_t)$ (and independently from the past given X_t).

Goal: Maximize (in expectation) the cumulative rewards.
Equivalently we want to minimize the cumulative regret

$$R_n = \mathbb{E} \sum_{t=1}^n \left(\max_{x \in \mathcal{X}} f(x) - Y_t \right).$$

\mathcal{X} -armed bandit game

Parameters available to the forecaster: the number of rounds n and the set of arms \mathcal{X} .

Parameters unknown to the forecaster: mean-payoff function $f : \mathcal{X} \rightarrow [0, 1]$, reward distributions (over $[0, 1]$) $M(x)$ such that $f(x)$ is the expectation of $M(x)$.

For each round $t = 1, 2, \dots, n$;

- 1 The player chooses an arm $X_t \in \mathcal{X}$.
- 2 The environment draws the reward Y_t from $M(X_t)$ (and independently from the past given X_t).

Goal: Maximize (in expectation) the cumulative rewards.
Equivalently we want to minimize the cumulative regret

$$R_n = \mathbb{E} \sum_{t=1}^n \left(\max_{x \in \mathcal{X}} f(x) - Y_t \right).$$

\mathcal{X} -armed bandit game

Parameters available to the forecaster: the number of rounds n and the set of arms \mathcal{X} .

Parameters unknown to the forecaster: mean-payoff function $f : \mathcal{X} \rightarrow [0, 1]$, reward distributions (over $[0, 1]$) $M(x)$ such that $f(x)$ is the expectation of $M(x)$.

For each round $t = 1, 2, \dots, n$;

- 1 The player chooses an arm $X_t \in \mathcal{X}$.
- 2 The environment draws the reward Y_t from $M(X_t)$ (and independently from the past given X_t).

Goal: Maximize (in expectation) the cumulative rewards.
Equivalently we want to minimize the cumulative regret

$$R_n = \mathbb{E} \sum_{t=1}^n \left(\max_{x \in \mathcal{X}} f(x) - Y_t \right).$$

\mathcal{X} -armed bandit game

Parameters available to the forecaster: the number of rounds n and the set of arms \mathcal{X} .

Parameters unknown to the forecaster: mean-payoff function $f : \mathcal{X} \rightarrow [0, 1]$, reward distributions (over $[0, 1]$) $M(x)$ such that $f(x)$ is the expectation of $M(x)$.

For each round $t = 1, 2, \dots, n$;

- ① The player chooses an arm $X_t \in \mathcal{X}$.
- ② The environment draws the reward Y_t from $M(X_t)$ (and independently from the past given X_t).

Goal: Maximize (in expectation) the cumulative rewards.
Equivalently we want to minimize the cumulative regret

$$R_n = \mathbb{E} \sum_{t=1}^n \left(\max_{x \in \mathcal{X}} f(x) - Y_t \right).$$

\mathcal{X} -armed bandit game

Parameters available to the forecaster: the number of rounds n and the set of arms \mathcal{X} .

Parameters unknown to the forecaster: mean-payoff function $f : \mathcal{X} \rightarrow [0, 1]$, reward distributions (over $[0, 1]$) $M(x)$ such that $f(x)$ is the expectation of $M(x)$.

For each round $t = 1, 2, \dots, n$;

- ① The player chooses an arm $X_t \in \mathcal{X}$.
- ② The environment draws the reward Y_t from $M(X_t)$ (and independently from the past given X_t).

Goal: Maximize (in expectation) the cumulative rewards.
Equivalently we want to minimize the cumulative regret

$$R_n = \mathbb{E} \sum_{t=1}^n \left(\max_{x \in \mathcal{X}} f(x) - Y_t \right).$$

Motivating examples

- **Calibrating the temperature** or levels of other inputs to a reaction so as to maximize the yield of a **chemical process**.
- **Pricing** a new product with **uncertain demand** in order to **maximize revenue**
- In general: online parameter tuning of numerical methods.
- Note: in the pricing problem **different product lines** could also be tested while tuning the price \Rightarrow **hybrid continuous/discrete set of arms**.

Motivating examples

- **Calibrating the temperature** or levels of other inputs to a reaction so as to maximize the yield of a **chemical process**.
- **Pricing** a new product with **uncertain demand** in order to **maximize revenue**
- In general: online parameter tuning of numerical methods.
- Note: in the pricing problem **different product lines** could also be tested while tuning the price \Rightarrow **hybrid continuous/discrete set of arms**.

Motivating examples

- **Calibrating the temperature** or levels of other inputs to a reaction so as to maximize the yield of a **chemical process**.
- **Pricing** a new product with **uncertain demand** in order to **maximize revenue**
- In general: online parameter tuning of numerical methods.
- Note: in the pricing problem **different product lines** could also be tested while tuning the price \Rightarrow **hybrid continuous/discrete set of arms**.

Motivating examples

- Calibrating the temperature or levels of other inputs to a reaction so as to maximize the yield of a chemical process.
- Pricing a new product with uncertain demand in order to maximize revenue
- In general: online parameter tuning of numerical methods.
- Note: in the pricing problem different product lines could also be tested while tuning the price \Rightarrow hybrid continuous/discrete set of arms.

Summary of the talk

- We present a new strategy, **Hierarchical Optimistic Optimization (HOO)**. It is based on a **tree-representation** of the search space, that we explore non-uniformly thanks to **upper confidence bounds** assigned to each nodes.
- Main theoretical result: if one knows the **local regularity** of the mean-payoff **function around its maximum**, then it is possible to obtain a cumulative regret of order \sqrt{n} .
- In particular, using n (noisy) evaluation of the function we can find **the maximum at a precision $1/\sqrt{n}$, independently of the ambient dimension!** Note that in a minimax sense, one can only find the maximum at a precision $n^{-1/(d+2)}$.

Summary of the talk

- We present a new strategy, **Hierarchical Optimistic Optimization (HOO)**. It is based on a **tree-representation** of the search space, that we explore non-uniformly thanks to **upper confidence bounds** assigned to each nodes.
- Main theoretical result: if one knows the **local regularity** of the mean-payoff **function around its maximum**, then it is possible to obtain a cumulative regret of order \sqrt{n} .
- In particular, using n (noisy) evaluation of the function we can find **the maximum at a precision $1/\sqrt{n}$, independently of the ambient dimension!** Note that in a minimax sense, one can only find the maximum at a precision $n^{-1/(d+2)}$.

Summary of the talk

- We present a new strategy, **Hierarchical Optimistic Optimization (HOO)**. It is based on a **tree-representation** of the search space, that we explore non-uniformly thanks to **upper confidence bounds** assigned to each nodes.
- Main theoretical result: if one knows the **local regularity** of the mean-payoff **function around its maximum**, then it is possible to obtain a cumulative regret of order \sqrt{n} .
- In particular, using n (noisy) evaluation of the function we can find **the maximum at a precision $1/\sqrt{n}$, independently of the ambient dimension!** Note that in a minimax sense, one can only find the maximum at a precision $n^{-1/(d+2)}$.

Local regularity around the maximum

Let ℓ be *dissimilarity* measure, that is, a non-negative mapping $\ell: \mathcal{X}^2 \rightarrow \mathbb{R}$ satisfying $\ell(x, x) = 0$.

Assumption (Weakly Lipschitz)

For all $x \in \mathcal{X}$ and $\epsilon \geq 0$, if $x \in \mathcal{X}_\epsilon = \{x \in \mathcal{X}, f^* - f(x) \leq \epsilon\}$ then for any $y \in \mathcal{X}$, $f(x) - f(y) \leq \max(\epsilon, \ell(x, y))$.

Local regularity around the maximum

Let ℓ be *dissimilarity* measure, that is, a non-negative mapping $\ell: \mathcal{X}^2 \rightarrow \mathbb{R}$ satisfying $\ell(x, x) = 0$.

Assumption (Weakly Lipschitz)

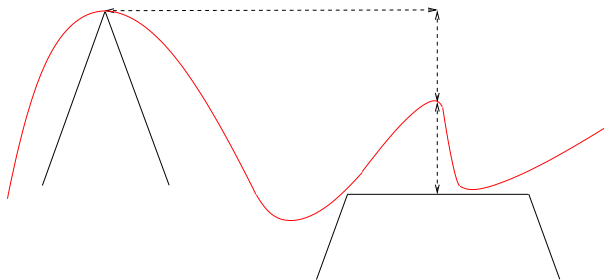
For all $x \in \mathcal{X}$ and $\epsilon \geq 0$, if $x \in \mathcal{X}_\epsilon = \{x \in \mathcal{X}, f^* - f(x) \leq \epsilon\}$ then for any $y \in \mathcal{X}$, $f(x) - f(y) \leq \max(\epsilon, \ell(x, y))$.

Local regularity around the maximum

Let ℓ be *dissimilarity* measure, that is, a non-negative mapping $\ell: \mathcal{X}^2 \rightarrow \mathbb{R}$ satisfying $\ell(x, x) = 0$.

Assumption (Weakly Lipschitz)

For all $x \in \mathcal{X}$ and $\epsilon \geq 0$, if $x \in \mathcal{X}_\epsilon = \{x \in \mathcal{X}, f^* - f(x) \leq \epsilon\}$ then for any $y \in \mathcal{X}$, $f(x) - f(y) \leq \max(\epsilon, \ell(x, y))$.



HOO - Input

- HOO receives as input a sequence $(\mathcal{P}_{h,i})_{h \geq 0, 1 \leq i \leq 2^h}$ of subsets of \mathcal{X} satisfying:
 - 1 $\mathcal{P}_{0,1} = \mathcal{X}$,
 - 2 $\mathcal{P}_{h,i} = \mathcal{P}_{h+1,2i-1} \cup \mathcal{P}_{h+1,2i}$.
 - 3 $\exists \rho \in (0, 1) : \text{diam}(\mathcal{P}_{h,i}) \leq \rho^h$ where
 $\text{diam}(\mathcal{P}_{h,i}) = \sup_{x,y \in \mathcal{P}_{h,i}} \ell(x,y)$.
- We view this as a tree where node (h,i) (at depth h and position i) is associated to the domain $\mathcal{P}_{h,i}$.

HOO - Input

- HOO receives as input a sequence $(\mathcal{P}_{h,i})_{h \geq 0, 1 \leq i \leq 2^h}$ of subsets of \mathcal{X} satisfying:
 - 1 $\mathcal{P}_{0,1} = \mathcal{X}$,
 - 2 $\mathcal{P}_{h,i} = \mathcal{P}_{h+1,2i-1} \cup \mathcal{P}_{h+1,2i}$.
 - 3 $\exists \rho \in (0, 1) : \text{diam}(\mathcal{P}_{h,i}) \leq \rho^h$ where
 $\text{diam}(\mathcal{P}_{h,i}) = \sup_{x,y \in \mathcal{P}_{h,i}} \ell(x,y)$.
- We view this as a tree where node (h,i) (at depth h and position i) is associated to the domain $\mathcal{P}_{h,i}$.

HOO - Input

- HOO receives as input a sequence $(\mathcal{P}_{h,i})_{h \geq 0, 1 \leq i \leq 2^h}$ of subsets of \mathcal{X} satisfying:
 - 1 $\mathcal{P}_{0,1} = \mathcal{X}$,
 - 2 $\mathcal{P}_{h,i} = \mathcal{P}_{h+1,2i-1} \cup \mathcal{P}_{h+1,2i}$.
 - 3 $\exists \rho \in (0, 1) : \text{diam}(\mathcal{P}_{h,i}) \leq \rho^h$ where
 $\text{diam}(\mathcal{P}_{h,i}) = \sup_{x,y \in \mathcal{P}_{h,i}} \ell(x,y)$.
- We view this as a tree where node (h,i) (at depth h and position i) is associated to the domain $\mathcal{P}_{h,i}$.

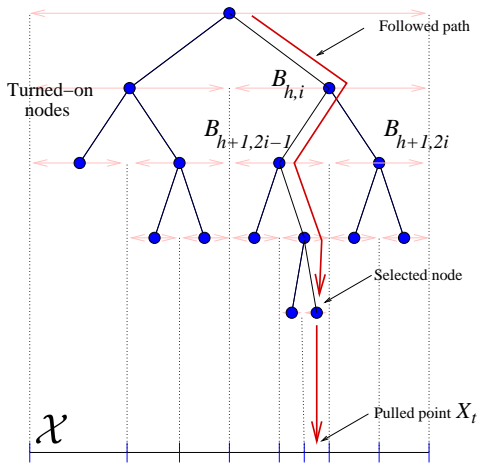
HOO - Input

- HOO receives as input a sequence $(\mathcal{P}_{h,i})_{h \geq 0, 1 \leq i \leq 2^h}$ of subsets of \mathcal{X} satisfying:
 - ① $\mathcal{P}_{0,1} = \mathcal{X}$,
 - ② $\mathcal{P}_{h,i} = \mathcal{P}_{h+1,2i-1} \cup \mathcal{P}_{h+1,2i}$.
 - ③ $\exists \rho \in (0, 1) : \text{diam}(\mathcal{P}_{h,i}) \leq \rho^h$ where
 $\text{diam}(\mathcal{P}_{h,i}) = \sup_{x,y \in \mathcal{P}_{h,i}} \ell(x, y)$.
- We view this as a tree where node (h, i) (at depth h and position i) is associated to the domain $\mathcal{P}_{h,i}$.

HOO - Input

- HOO receives as input a sequence $(\mathcal{P}_{h,i})_{h \geq 0, 1 \leq i \leq 2^h}$ of subsets of \mathcal{X} satisfying:
 - ① $\mathcal{P}_{0,1} = \mathcal{X}$,
 - ② $\mathcal{P}_{h,i} = \mathcal{P}_{h+1,2i-1} \cup \mathcal{P}_{h+1,2i}$.
 - ③ $\exists \rho \in (0, 1) : \text{diam}(\mathcal{P}_{h,i}) \leq \rho^h$ where
 $\text{diam}(\mathcal{P}_{h,i}) = \sup_{x,y \in \mathcal{P}_{h,i}} \ell(x, y)$.
- We view this as a tree where node (h, i) (at depth h and position i) is associated to the domain $\mathcal{P}_{h,i}$.

HOO - Global strategy given B -values for each node



HOO - Definition of B -values

- Let $T_{h,i}(n)$ be the number of points we pulled in (h, i) .
- Let $\hat{\mu}_{h,i}(n)$ be the empirical average in the domain (h, i) .
- We consider the following upper confidence bound for each node already visited :

$$U_{h,i}(n) = \hat{\mu}_{h,i}(n) + \sqrt{\frac{2 \ln n}{T_{h,i}(n)}} + \text{diam}(\mathcal{P}_{h,i}).$$

- Our B -values are defined as:

$$B_{h,i}(n) = \min \left\{ U_{h,i}(n), \max \{ B_{h+1,2i-1}(n), B_{h+1,2i}(n) \} \right\}.$$

HOO - Definition of B -values

- Let $T_{h,i}(n)$ be the number of points we pulled in (h, i) .
- Let $\hat{\mu}_{h,i}(n)$ be the empirical average in the domain (h, i) .
- We consider the following upper confidence bound for each node already visited :

$$U_{h,i}(n) = \hat{\mu}_{h,i}(n) + \sqrt{\frac{2 \ln n}{T_{h,i}(n)}} + \text{diam}(\mathcal{P}_{h,i}).$$

- Our B -values are defined as:

$$B_{h,i}(n) = \min \left\{ U_{h,i}(n), \max \{ B_{h+1,2i-1}(n), B_{h+1,2i}(n) \} \right\}.$$

HOO - Definition of B -values

- Let $T_{h,i}(n)$ be the number of points we pulled in (h, i) .
- Let $\hat{\mu}_{h,i}(n)$ be the empirical average in the domain (h, i) .
- We consider the following upper confidence bound for each node already visited :

$$U_{h,i}(n) = \hat{\mu}_{h,i}(n) + \sqrt{\frac{2 \ln n}{T_{h,i}(n)}} + \text{diam}(\mathcal{P}_{h,i}).$$

- Our B -values are defined as:

$$B_{h,i}(n) = \min \left\{ U_{h,i}(n), \max \{ B_{h+1,2i-1}(n), B_{h+1,2i}(n) \} \right\}.$$

HOO - Definition of B -values

- Let $T_{h,i}(n)$ be the number of points we pulled in (h, i) .
- Let $\hat{\mu}_{h,i}(n)$ be the empirical average in the domain (h, i) .
- We consider the following upper confidence bound for each node already visited :

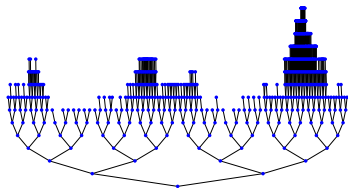
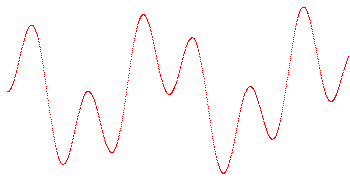
$$U_{h,i}(n) = \hat{\mu}_{h,i}(n) + \sqrt{\frac{2 \ln n}{T_{h,i}(n)}} + \text{diam}(\mathcal{P}_{h,i}).$$

- Our B -values are defined as:

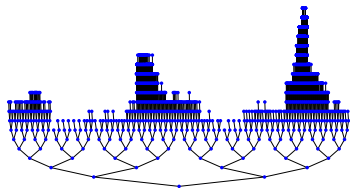
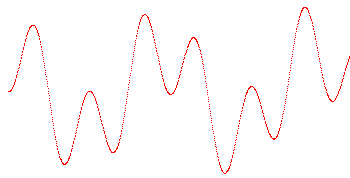
$$B_{h,i}(n) = \min \left\{ U_{h,i}(n), \max \{ B_{h+1,2i-1}(n), B_{h+1,2i}(n) \} \right\}.$$

HOO - Numerical Example

$n = 1000$



$n = 10000$



Main result

Definition (Near-optimality dimension)

Let $d \geq 0$ be such that $\mathcal{X}_\epsilon = \{x \in \mathcal{X}, f^* - f(x) \leq \epsilon\}$ can be packed with $O(\epsilon^{-d})$ balls of radius ϵ .

Main result

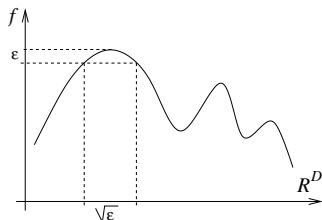
Definition (Near-optimality dimension)

Let $d \geq 0$ be such that $\mathcal{X}_\epsilon = \{x \in \mathcal{X}, f^* - f(x) \leq \epsilon\}$ can be packed with $O(\epsilon^{-d})$ balls of radius ϵ .

Main result

Definition (Near-optimality dimension)

Let $d \geq 0$ be such that $\mathcal{X}_\epsilon = \{x \in \mathcal{X}, f^* - f(x) \leq \epsilon\}$ can be packed with $O(\epsilon^{-d})$ balls of radius ϵ .



- $\ell(x, y) = \|x - y\| \Rightarrow d = D/2.$
- $\ell(x, y) = \|x - y\|^2 \Rightarrow d = 0.$

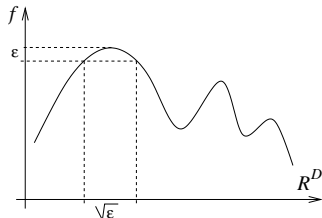
Theorem

HOO satisfy $R_n \leq \tilde{O}(n^{(d+1)/(d+2)}).$

Main result

Definition (Near-optimality dimension)

Let $d \geq 0$ be such that $\mathcal{X}_\epsilon = \{x \in \mathcal{X}, f^* - f(x) \leq \epsilon\}$ can be packed with $O(\epsilon^{-d})$ balls of radius ϵ .



- $\ell(x, y) = \|x - y\| \Rightarrow d = D/2$.
- $\ell(x, y) = \|x - y\|^2 \Rightarrow d = 0$.

Theorem

HOO satisfy $R_n \leq \tilde{O}(n^{(d+1)/(d+2)})$.

Example

$\mathcal{X} = [0, 1]^D$, $\alpha \geq 0$ and f locally " α -smooth" around (any of) its maximum x^* (in finite number):

$$f(x^*) - f(x) = \Theta(\|x - x^*\|^\alpha) \text{ as } x \rightarrow x^*.$$

Theorem

Assume that we run HOO with diameters measured with $\ell(x, y) = \|x - y\|^\beta$.

- **Known smoothness:** $\beta = \alpha$. $R_n \leq \tilde{O}(\sqrt{n})$, i.e., the rate is independent of the dimension D . Previously known for $D = 1$ or $\alpha \leq 1$.
- **Smoothness underestimated:** $\beta < \alpha$.
 $R_n \leq \tilde{O}(n^{(d+1)/(d+2)})$ where $d = D \left(\frac{1}{\beta} - \frac{1}{\alpha} \right)$.
- **Smoothness overestimated:** $\beta > \alpha$. No guarantee. Note: UCT corresponds to $\beta = +\infty$.

Example

$\mathcal{X} = [0, 1]^D$, $\alpha \geq 0$ and f locally " α -smooth" around (any of) its maximum x^* (in finite number):

$$f(x^*) - f(x) = \Theta(\|x - x^*\|^\alpha) \text{ as } x \rightarrow x^*.$$

Theorem

Assume that we run HOO with diameters measured with $\ell(x, y) = \|x - y\|^\beta$.

- **Known smoothness:** $\beta = \alpha$. $R_n \leq \tilde{O}(\sqrt{n})$, i.e., the rate is independent of the dimension D . Previously known for $D = 1$ or $\alpha \leq 1$.
- **Smoothness underestimated:** $\beta < \alpha$.
 $R_n \leq \tilde{O}(n^{(d+1)/(d+2)})$ where $d = D \left(\frac{1}{\beta} - \frac{1}{\alpha} \right)$.
- **Smoothness overestimated:** $\beta > \alpha$. No guarantee. Note: UCT corresponds to $\beta = +\infty$.

Example

$\mathcal{X} = [0, 1]^D$, $\alpha \geq 0$ and f locally " α -smooth" around (any of) its maximum x^* (in finite number):

$$f(x^*) - f(x) = \Theta(\|x - x^*\|^\alpha) \text{ as } x \rightarrow x^*.$$

Theorem

Assume that we run HOO with diameters measured with $\ell(x, y) = \|x - y\|^\beta$.

- **Known smoothness:** $\beta = \alpha$. $R_n \leq \tilde{O}(\sqrt{n})$, i.e., the rate is independent of the dimension D . Previously known for $D = 1$ or $\alpha \leq 1$.
- **Smoothness underestimated:** $\beta < \alpha$.
 $R_n \leq \tilde{O}(n^{(d+1)/(d+2)})$ where $d = D \left(\frac{1}{\beta} - \frac{1}{\alpha} \right)$.
- **Smoothness overestimated:** $\beta > \alpha$. No guarantee. Note: UCT corresponds to $\beta = +\infty$.

Example

$\mathcal{X} = [0, 1]^D$, $\alpha \geq 0$ and f locally " α -smooth" around (any of) its maximum x^* (in finite number):

$$f(x^*) - f(x) = \Theta(\|x - x^*\|^\alpha) \text{ as } x \rightarrow x^*.$$

Theorem

Assume that we run HOO with diameters measured with $\ell(x, y) = \|x - y\|^\beta$.

- **Known smoothness:** $\beta = \alpha$. $R_n \leq \tilde{O}(\sqrt{n})$, i.e., the rate is independent of the dimension D . Previously known for $D = 1$ or $\alpha \leq 1$.
- **Smoothness underestimated:** $\beta < \alpha$.
 $R_n \leq \tilde{O}(n^{(d+1)/(d+2)})$ where $d = D \left(\frac{1}{\beta} - \frac{1}{\alpha} \right)$.
- **Smoothness overestimated:** $\beta > \alpha$. No guarantee. Note: UCT corresponds to $\beta = +\infty$.

Example

$\mathcal{X} = [0, 1]^D$, $\alpha \geq 0$ and f locally " α -smooth" around (any of) its maximum x^* (in finite number):

$$f(x^*) - f(x) = \Theta(\|x - x^*\|^\alpha) \text{ as } x \rightarrow x^*.$$

Theorem

Assume that we run HOO with diameters measured with $\ell(x, y) = \|x - y\|^\beta$.

- **Known smoothness:** $\beta = \alpha$. $R_n \leq \tilde{O}(\sqrt{n})$, i.e., the rate is independent of the dimension D . Previously known for $D = 1$ or $\alpha \leq 1$.
- **Smoothness underestimated:** $\beta < \alpha$.
 $R_n \leq \tilde{O}(n^{(d+1)/(d+2)})$ where $d = D \left(\frac{1}{\beta} - \frac{1}{\alpha} \right)$.
- **Smoothness overestimated:** $\beta > \alpha$. No guarantee. Note: UCT corresponds to $\beta = +\infty$.